

# First-Order Interpolation Derived from Propositional Interpolation<sup>\*</sup>

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Ever since Craig's seminal paper on interpolation [3], interpolation properties have been recognized as important properties of logical systems. Recall that a logic  $L$  has *interpolation* if whenever  $A \rightarrow B$  holds in  $L$  there exists a formula  $I$  in the common language of  $A$  and  $B$  such that  $A \rightarrow I$  and  $I \rightarrow B$  hold in  $L$ .

Propositional interpolation properties can be determined and classified with relative ease using the ground-breaking results of Maksimova cf. [7, 6, 5]. This approach is based on an algebraic analysis of the logic in question. In contrast first-order interpolation properties are notoriously hard to determine, even for logics where propositional interpolation is more or less obvious. For example it is unknown whether  $G_{[0,1]}^{\text{QF}}$  (first-order infinitely-valued Gödel logic) interpolates (cf. [1]) and even for  $\text{MC}^{\text{QF}}$ , the logic of constant domain Kripke frames of three worlds with two top worlds (an extension of MC), interpolation proofs are very hard cf. Ono [8]. This situation is due to the lack of an adequate algebraization of non-classical first-order logics. In this paper we present a proof theoretic methodology to reduce first-order interpolation to propositional interpolation:

$$\left. \begin{array}{l} \text{existence of suitable skolemizations} + \\ \text{existence of Herbrand expansions} + \\ \text{propositional interpolation} \end{array} \right\} \Rightarrow \begin{array}{l} \text{first-order} \\ \text{interpolation.} \end{array}$$

The construction of the first-order interpolant from the propositional interpolant follows this procedure:

1. Develop a validity equivalent skolemization replacing all strong quantifiers<sup>3</sup> in the valid formula  $A \rightarrow B$  to obtain the valid formula  $A_1 \rightarrow B_1$ .
2. Construct a valid Herbrand expansion  $A_2 \rightarrow B_2$  for  $A_1 \rightarrow B_1$ . Occurrences of  $\exists xB(x)$  and  $\forall xA(x)$  are replaced by suitable finite disjunctions  $\bigvee B(t_i)$  and conjunctions  $\bigwedge B(t_i)$ , respectively.
3. Interpolate the propositionally valid formula  $A_2 \rightarrow B_2$  with the propositional interpolant  $I^*$ :  $A_2 \rightarrow I^*$  and  $I^* \rightarrow B_2$  are propositionally valid.
4. Reintroduce weak quantifiers to obtain valid formulas  $A_1 \rightarrow I^*$  and  $I^* \rightarrow B_1$ .

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<sup>\*</sup> This abstract is based on the publication [2].

<sup>3</sup> Here we are dealing with quantifiers  $\forall$  and  $\exists$  such that  $A(t) \rightarrow \exists xA(x)$  and  $\forall xA(x) \rightarrow A(t)$  hold. This occurrence of quantifiers is called weak, the dual occurrence is called strong.

5. Eliminate all function symbols and constants not in the common language of  $A_1$  and  $B_1$  by introducing suitable quantifiers in  $I^*$  (note that no Skolem functions are in the common language, therefore they are eliminated). Let  $I$  be the result.
6.  $I$  is an interpolant for  $A_1 \rightarrow B_1$ .  $A_1 \rightarrow I$  and  $I \rightarrow B_1$  are skolemizations of  $A \rightarrow I$  and  $I \rightarrow B$ . Therefore  $I$  is an interpolant of  $A \rightarrow B$ .

It is decidable if propositional lattice based finitely-values logics admit the interpolation property [2]. Consequently, it is decidable if finitely-valued first-order logics admit the interpolation property. In this lecture we extend the methodology to prenex fragments of non-classical logics where Skolemization is admissible due to the second epsilon theorem [4].

## References

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